

# Uncommon Sense?

A smorgasbord of Coherence  
production advice,  
quizzes, & miscellany

(David Whitmarsh and Phil Wheeler Credit Suisse)

Aimed at those with low or intermediate levels of experience with Coherence, but even expert users may learn something new.

# Start Easy

# Production Clusters

What's your cluster's name?

# Production Clusters

What's your mode?

QUIZ!

# What happens?

- a. `cache.keySet().remove(key);`
- b. `cache.keySet(filter).remove(key);`

# JVM settings

# Set initial and max JVM sizes

-Xms3G –Xmx3G

Allocate all physical memory upfront

No nasty surprises as data volumes grow and physical memory becomes exhausted.

# Die on OOM

JVM behaviour is unpredictable and may endanger the cluster once OOM has occurred.

## UNIX

`-XX:OnOutOfMemoryError="kill -9 %p"`

## Windows

`-XX:OnOutOfMemoryError="taskkill /F /PID %p"`

# Heap dump on OOM

Heap dumps assist your diagnosis:

`-XX:+HeapDumpOnOutOfMemoryError`

`-XX:HeapDumpPath=<node-specific-file>`

QUIZ!

# Lively up your cache?

```
public class BackingMapListener implements MapListener {  
  
    @Override  
    public void entryInserted(MapEvent e) {  
        // spend 5s doing some work  
    }  
    ...  
}
```

How many puts can be handled?

Configure 10 threads: Now how many?

# Garbage Collection

# Concurrent Mark Sweep

Short full GC times for relatively little overhead

-XX:+UseConcMarkSweepGC

-XX:CMSInitiatingOccupancyFraction=75

-XX:+UseInitiatingOccupancyOnly

**How do you decide the correct occupancy fraction?**

- Tie this to memory threshold alerts.
- Heap appears to grow until this level is hit, initiating a full GC.
- Alerting below this level will cause false alarms.

# Log Garbage Collections

-Xloggc:<node-specific-file>

-XX:+PrintGCDetails

For more verbose logging:

-XX:+PrintGCDetails

-XX:+PrintTenuringDistribution

# JVM sizing & capacity planning

# How much heap for your data?

$$M = \frac{\sum_{i=0}^n \left( \frac{2s_i + x_i}{p(H - 1)} \right) c_i + \sum_{j=0}^m c_j s_j + w + h}{w}$$

M = total heap per storage node

C<sub>i</sub> = count of objects in distributed cache i

s<sub>i</sub> = size of objects in distributed cache i

x<sub>i</sub> = size of indexes for an entry in distributed cache i

H = number of hosts

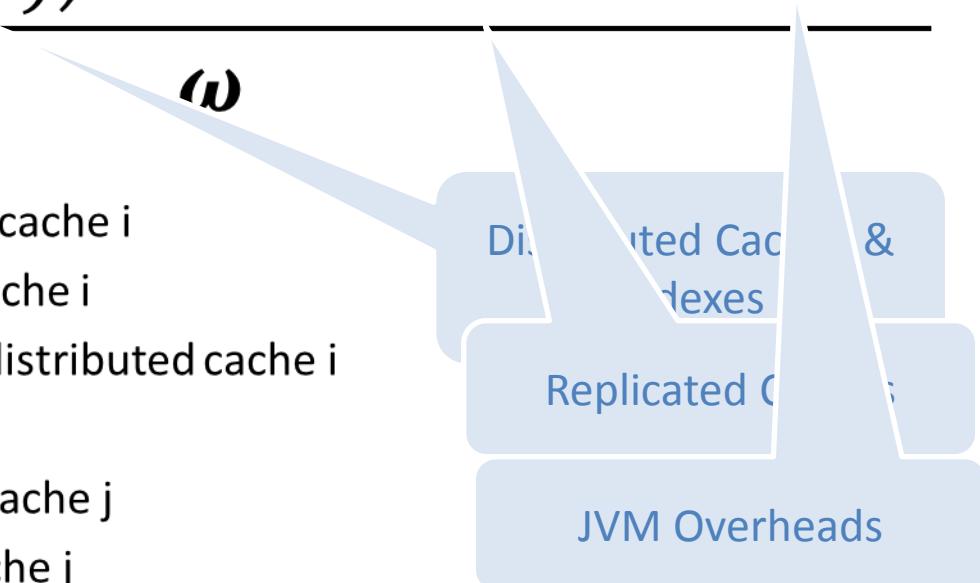
c<sub>j</sub> = count of objects in replicated cache j

s<sub>j</sub> = size of objects in replicated cache j

p = number of storage node processes per host

w = working heap overhead per jvm

h = fixed heap overhead per jvm



# How much physical memory do you need for your JVM?

$$V = a + P + Hb$$

V = OS virtual image size

H = Heap size

P = Permgen size

a, b are constants

For 32bit: a = 130MB, b = 1.07

For 64bit: a = 285MB, b = 1.05

From tests on 1.6 Hotspot JVM

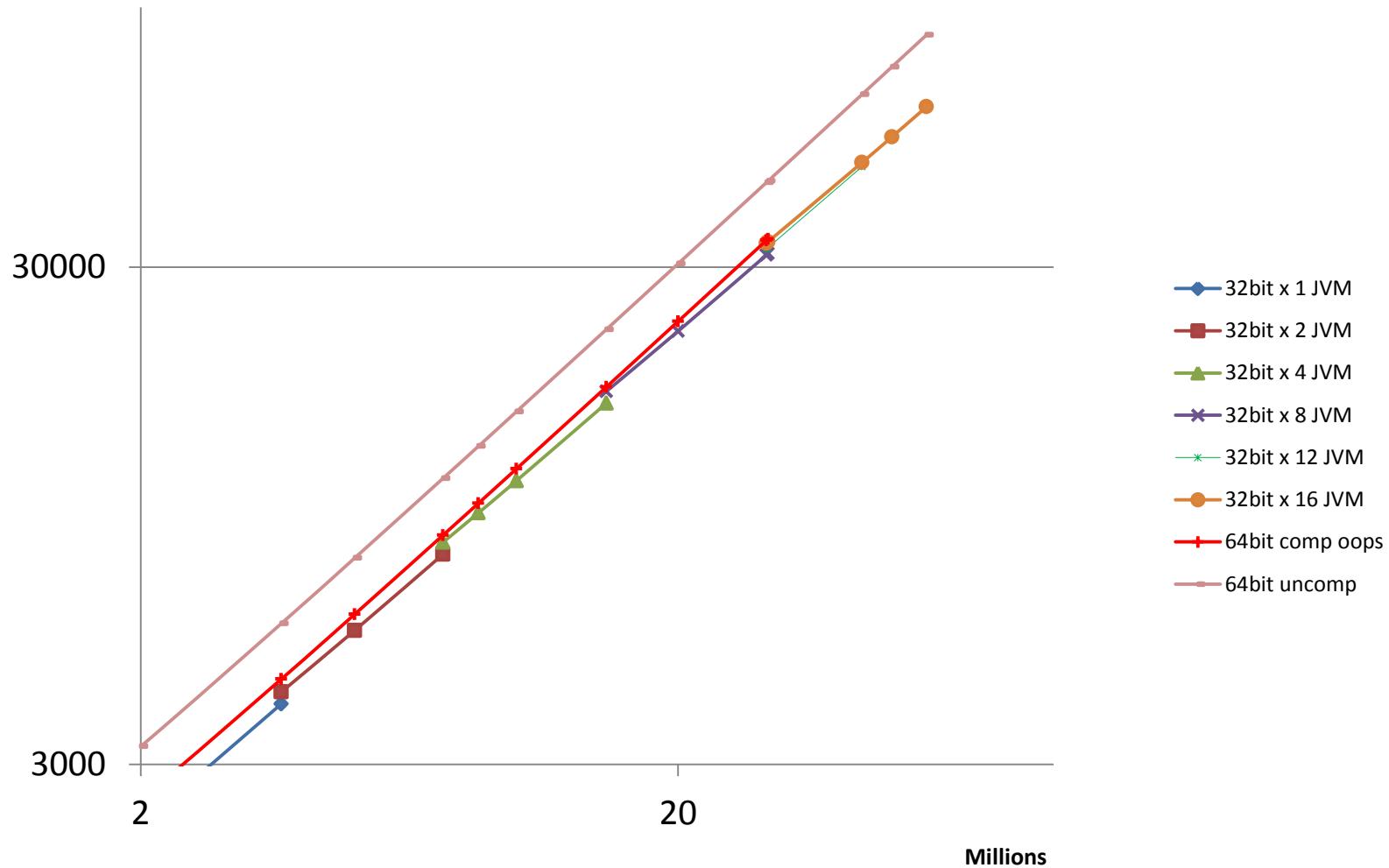
# Fewer, larger JVMs or more, smaller JVMs?

- GC times
- Efficiency of memory use
- Repartitioning time
- Write-behind bottleneck
- Filter query latency
- Result set sizes
- CPU Cores

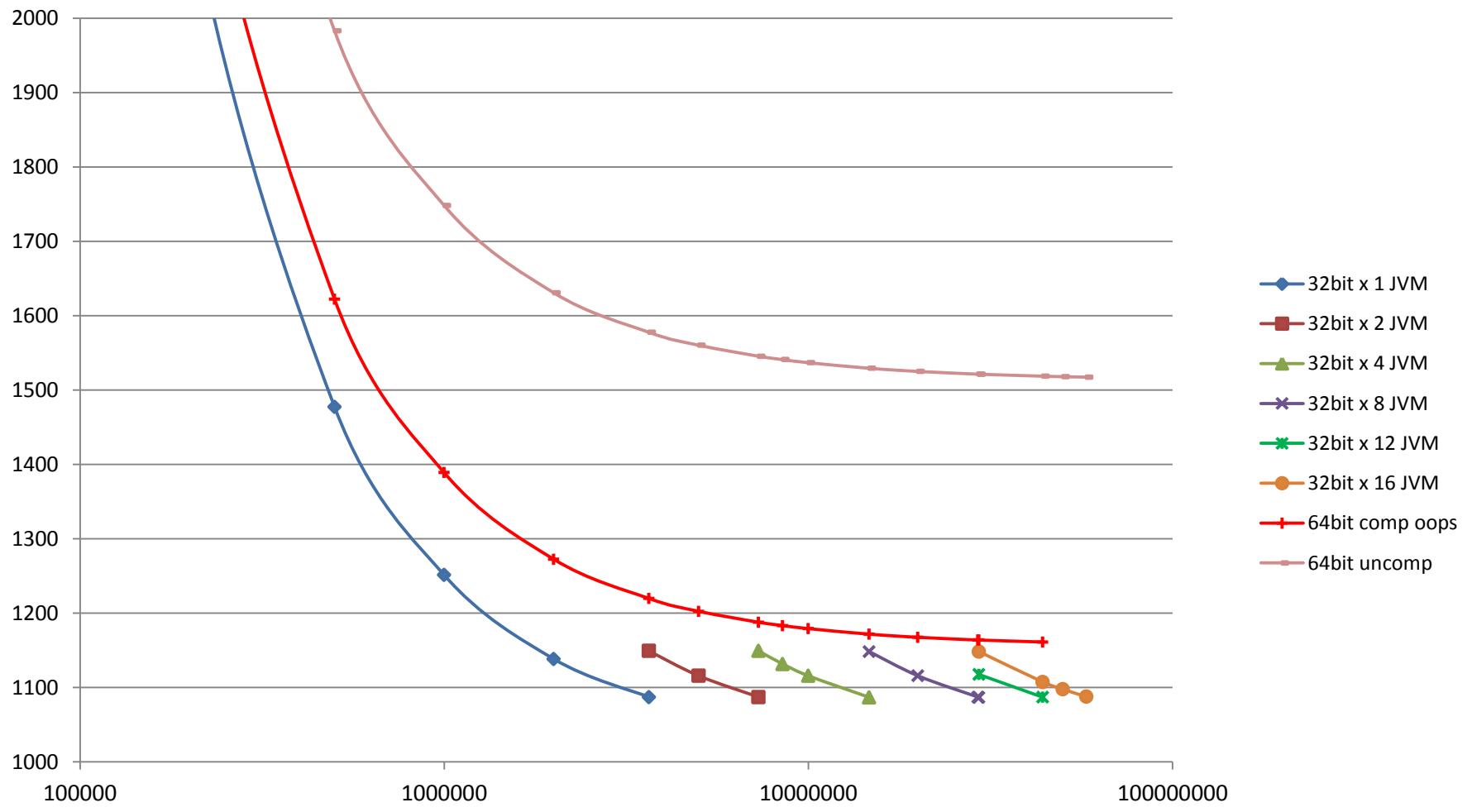
# 32bit vs. 64bit JVMs

- Heap limit **32bit**  $\approx$  **3.5GB**
- Heap limit **64bit** with compressedOops is **32GB**
- Throughput performance is comparable
- **64bit** without compressedOops – you don't want to do that!
- Objects take more space in **64bit**
  - +10% even with compressed oops
  - +40% with uncompressed

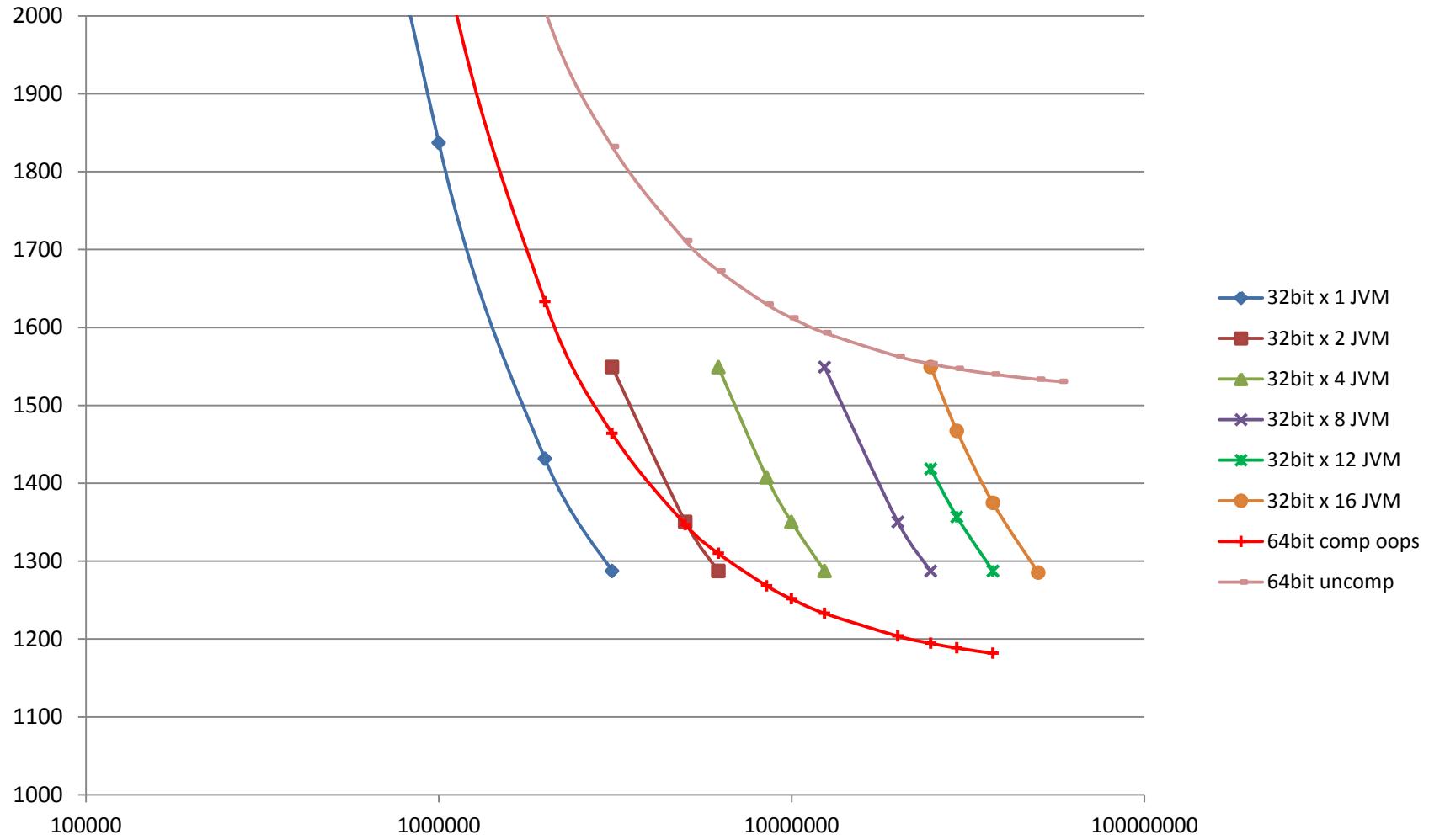
# Total Memory vs. Entries



# Memory Per Entry (low o/h)

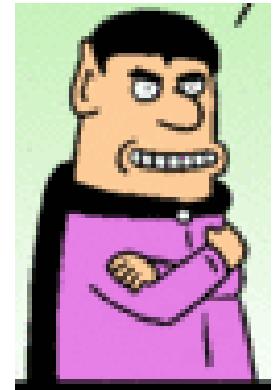


# Memory per Entry (high o/h)



# How much uncommitted physical memory do you need on a server?

- **Swapping** is extremely bad news.
- The most likely time to notice pages swapped out is during GC
- Some UNIX sysadmins standardise backup or monitoring tools that gobble up **huge** amounts of memory when they run.



# Can we prevent swapping?

- Leave lots of OS memory uncommitted.  
How much?
- Swappiness
- Huge pages
- sudo swapoff –a
- mlockall



# RAM Pinning with `mlockall`

- `mlockall` : standard C library function
- Locks a process's virtual address space into physical memory
- Requires the memlock user limit to be set in `/etc/security/limits.conf`
- Can be called via JNA

# RAM Pinning with mlockall

```
package com.csg.dtacc.coherence.utils;

import com.sun.jna.Library;
import com.sun.jna.Native;

public class MemLock {

    public static final int MCL_CURRENT = 1;
    public static final int MCL_FUTURE = 2;

    private interface CLibrary extends Library {
        int mlockall(int flags);
    }

    private synchronized static void mlockall() {
        CLibrary instance = (CLibrary) Native.loadLibrary("c", CLibrary.class);
        return instance.mlockall(MCL_CURRENT | MCL_FUTURE);
    }
}
```

# Checking the Results

Check result by looking at `/proc/$pid/status`

```
Name:      java
State:    S (sleeping)
...
VmPeak:   3960000 kB
VmSize:   3959996 kB
VmLck:    3959996 kB
VmHWM:    3958052 kB
VmRSS:    3958048 kB
VmData:   3945308 kB
VmStk:     40 kB
VmExe:    40 kB
VmLib:   14136 kB
VmPTE:   7752 kB
Threads:   47
...
```

**VmLck** : shows the amount of memory locked for the process

# Miscellany

QUIZ!

# What happens when an Exception is thrown?

- a. in entryProcessor.process()

```
Obj result = cache.invoke(key, entryProcessor);
```

*In 3.7.1.0 ?*

- b. in entryProcessor.process()

```
Map result = cache.invokeAll(keyset, entryProcessor);
```

- c. in filter.evaluate()

```
Set result = cache.entrySet(filter);
```

- d. in entryProcessor.process() or filter.evaluate()

```
Map result = cache.invokeAll(filter, entryProcessor);
```

# POF testing

```
public class ValueObjectTest {  
  
    private static final String POF_CONFIG_XML = "pof-config.xml";  
  
    @Test  
    public void testPofFidelity() {  
        ValueObject vo = new ValueObject("constructor", "args");  
        assertPofFidelity(vo);  
    }  
  
    private void assertPofFidelity(Object example) {  
        ConfigurablePofContext ctxt = new ConfigurablePofContext(POF_CONFIG_XML);  
  
        Binary binary = ExternalizableHelper.toBinary(example, ctxt);  
        Object result = ExternalizableHelper.fromBinary(binary, ctxt);  
  
        assertEquals(example, result);  
    }  
}
```

# POF testing using commons-lang EqualsBuilder

```
private void assertPofFidelity(Object example) {
    ConfigurablePofContext ctxt = new ConfigurablePofContext(POF_CONFIG_XML);

    Binary binary = ExternalizableHelper.toBinary(example, ctxt);
    Object result = ExternalizableHelper.fromBinary(binary, ctxt);

    assertTrue(EqualsBuilder.reflectionEquals(example, result));
}

}
```

# POF Testing in .NET

```
[TestFixture]
public class POFTests
{
    ...
    [Test]
    public void TestPofFidelity()
    {
        ConfigurablePofContext ctx = new ConfigurablePofContext("pof-config.xml");

        Object original = new ValueObject("constructor", "args");
        Binary bin   = SerializationHelper.ToBinary(original, ctx);
        Object copy  = SerializationHelper.FromBinary(bin, ctx);

        Assert.AreEqual(original, copy);
    }
}
```

# Don't start replicated service

- Replicated data is copied to every node. Where do you use it?
  - EntryProcessors in storage nodes?
  - Application logic in storage-disabled nodes?
  - Proxy nodes?
  - JMX nodes? Really?
- Start it only where needed
  - With a **system-property**
  - Lazy start on access to a replicated cache

```
<replicated-scheme>
    <scheme-name>example-replicated</scheme-name>
    ...
    <autostart system-property="replicated.service.enabled">false</autostart>
</replicated-scheme>
```

QUIZ!

# Threads

**my-distributed-service has**

<thread-count>**5**</thread-count>,

**...is configured for write-behind,**

**...and has 2 caches.**

**How many *connections* do I need in my JDBC pool so that no thread will ever wait for a connection?**